Derived completeness



Classica/I-**Classical completions behaving badly** A=nrg I finitelyse nexted iteal McMod N's christaly I-complete if M-) I'm MIM

(m) my isinjective = M's Indically squated) - classificaty I complete modules we not an atelian content of the exact (in the some)

- MI - M 15 nt ever of the exact

(by to the presure fathers.)

- complete does not presure fathers. it Interse, her A & Minsuell

An example of Yekutieli A = R I = RM2= { (XnE Min Xn/n -) 0} Z/ PT)

The milds Mo, M, Mz is complete, no/m, M//nz. ut at Mo/my: dosmed Mz, n Mo is Mi. 0 7 M2 -1 M, -> Mo/M2 -1 0 completion takes 0 -> M2 -> M0/M, -> 0

Derived I-complete modules

T = Inse Mez 1 MEMORA IS delived I-aylete it YftI, Homa (Af, M) = Ext'A (Af, M)=0 Note: sne O JACT) JAF DO IS a projective a sinker, X(I-TF) T D+ MIS SPINAVOLUTE TO ALM) TO KNOW. note: this walthing (#I har given t = 5 (#I tow eng multiple of f.

More on the definition of derived completeness Cenn for sine M, tresitot FEA to mhich (H) hulds is handical ideal Jut A. Pt: fcI, 16 A => fg (I). - f(I(=) f ~ EI (4 = 4m) - f. g EI 0-) A F+5 > A F(A3) DA g(A3) -> AF3(A3) -> 0 +(-5) = Mun Holen = 1 F3 G I, There to check (1) m's to set of sweeters

of I, or some ideal who swe radical:

<u>Classical vs. derived completeness</u> Lenna: 1) M Massically I amplete = M decient I applete

2) if M derven I - ramplete = M => lim M, fr.m.

=) classically I - complete = deriven I - imprete

+ I - a d, ally reported.

1) Homa (A+, M) = Homa (A+, lim M/Inm) = 0

V+FI

E + A (A+, M) = 0 0 -> M -> E -> A+ >> 0

E + A (A+, M) = 0 0 -> M -> E -> A+ >> 0 ene E mypons Wt neAt, In= Fent = eneM Mcmplete=> (mlehre en=h+ontflons+ f25/n12+... =) fenn=en, => segvence splits.

Classical vs. derived completeness

2) Che M M M/M (exerue: fins h

have) Give Xo, Xi... CM = XXXXXXXXXXXXX $E = Me(P) A e M(x_n - fen + ten) - 2 A f$ $= Me(P) A e M(x_n - fen + ten) - 3 f - n$ oute: M/fmM=E/fm = (snahe lema) a) M-1=-1Ax-10 Extenson, 11+3/y hypothesis, 10 7 xm 4. X+lo = x-xv+fe, = x-xv+X, +fez.

The category of derived complete modules A = 0.05DNP NAK 1) (Derved Nakayana) M = derved I-unpleke then M=D (=) M/IM = 0 2) in Obser for reved I-rays lete 4 middles to Mod A has a left adjust, de wed I-rapkhon The fill sheaterngot Mod a consistens of chimed I-complete models of abelian, and formation of herre's & where's warmtes with indusion. The category of derived-complete modules

(or if A is desired I - worlete -) I = Rad (A) (apply derived -) eny Kn, Hy sorsett A-modhis Level I cample to. Coment: , reste example, Mo/Mz is not I -adrially represented by the solution of the so

Derived completion

I=(h...) for)

derived I-royale for for chiris

composition of derived fi-royale for

for i=1.....

Derived completion: principal case lemail derold frage trans for Deris (WA -)Af -> At/A -) V)

M = Ext (AFA M.) 2) The vanathal exact sequence 0-) R'ImM(f)-) -> 1, m M/n m > 0 · Lustin my, M(f^n) ~M(f^n), s xf.

Cases where derived completions are classical Umma A=nns, I=(f) pnulpalidect

M=A-mod, If M has bounded f -hosion (i1 _ M(f00)= (/M(f1) = swe M(f1)) the derived I - ingliture danical I-rapleton. Copolly: 12=petectornsof Chr. to my f(P), R(f)=R(f), so

R has bunded f = -two m as a module over itself.

Photograph of the property of the property of the constrainty of the property of

Derived completeness/completion in a derived category

Similarly for KCD(A) demed antony at Mod A Kis derived I - cop lete of Rising (AF, K)=8 (My, to Medingenetis) YFEI. still Lilds: Derre 1 Nahay wa Krever I. wylet, KB 5 A, I=0 in D(A/I) => K=V , ~ D(A).

· Neved I-wyhte db, ich, + D(A) forma fall triang Sated show tegory uma left adjunt/level I-carpéhin) o Objet Ky Nevel I-inglete (5) 17'(K) are derved Implet MEA, Inved I-wy When as a modele 15 MC devel I-ay behin of MCOD) WA there yet be one terms.