Recent progress in computing zeta functions of varieties

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Based in part on joint work (in preparation) with Edgar Costa (NYU/ICERM/Dartmouth) and David Harvey (U. New South Wales).

Zeta functions of algebraic varieties

Let \mathbb{F}_q be a finite field of characteristic *p*. After Artin, Schmidt, and Weil, we define the *zeta function* of a variety *X* over \mathbb{F}_q as the formal Dirichlet series

$$\zeta(X,s) = \prod_{x} (1 - \#\kappa(x)^{-s})^{-1},$$

where x runs over closed points of X and $\kappa(x)$ denotes the residue field. (Equivalently, x runs over Galois orbits of $\overline{\mathbb{F}}_{q}$ -rational points and $\kappa(x)$ denotes the minimal field of definition.)

From now on, we write ζ as a formal power series in $T = q^{-s}$. Then

$$\zeta(X,T) = \exp\left(\sum_{n=1}^{\infty} \frac{T^n}{n} \# X(\mathbb{F}_{q^n})\right).$$

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Examples of zeta functions

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one can compute $\zeta(X, T)$ in some explicit examples. For one:

$$\zeta(\mathbb{P}^d_{\mathbb{F}_q}, \mathcal{T}) = rac{1}{(1-\mathcal{T})(1-q\mathcal{T})\cdots(1-q^d\mathcal{T})}$$

For another, if X is an elliptic curve over \mathbb{F}_q , then

$$\zeta(X,T) = \frac{1 - aT + qT^2}{(1 - T)(1 - qT)}, \qquad a = q + 1 - \#X(\mathbb{F}_q).$$

Based on these (and more) examples, Weil predicted that $\zeta(X, T)$ obeys analogues of the properties of the Riemann zeta function (analytic continuation, functional equation, Riemann hypothesis).

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The first of the Weil conjectures on zeta functions of algebraic varieties is:

Theorem

The power series $\zeta(X, T)$ represents a rational function of T.

This is widely known as a consequence of the construction of *étale cohomology* by Grothendieck et al. However, that was not the first proof!

Theorem (Dwork, 1960)

The power series $\zeta(X, T)$ is p-adic meromorphic: it is the ratio of two power series over \mathbb{Q}_p with infinite radii of convergence.

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We are interested in algorithms that given an explicit definition of X (i.e., defining equations), returns the rational function $\zeta(X, T)$. Why?

- For X a curve over \mathbb{F}_q , $\zeta(X, T)$ determines $\#X(\mathbb{F}_q)$. Questions about this count, especially its extreme values, were originally inspired by algebraic coding theory (Goppa construction) and have been studied for decades; see http://manypoints.org.
- For X a curve over \mathbb{F}_q , $\zeta(X, T)$ determines $\#J(X)(\mathbb{F}_q)$. This is relevant for cryptography, especially when X is of genus ≤ 3 .
- For X a variety over F_q, ζ(X, T) contains (via the Tate conjecture) information about algebraic cycles on X, e.g., the Picard number.
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- Special values of these L-functions carry arithmetic information via many conjectures (Birch–Swinnerton-Dyer, Beilinson, Deligne, etc.).
- Additional arithmetic information can be seen in the statistics of variation of Euler factors, via many conjectures (Sato-Tate, Fité-Rotger-K-Sutherland, Serre, etc.).
- These L-functions are conjecturally related to automorphic forms via the Langlands correspondence. One hopes both to collect evidence to help make precise new cases of the correspondence (e.g., genus 2 curves and Siegel modular forms), and to use existing results to compute automorphic L-functions.
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This is almost never practical! One can occasionally gain something by exploiting geometric information, e.g., automorphisms or correspondences which impose constraints on the factorization of $\zeta(X, T)$.

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Zeta functions via Weil cohomology

A Weil cohomology theory provides an interpretation

$$\zeta(X,T)=\prod_i \det(1-TF_i,V_i)^{(-1)^{i+1}}$$

for some endomorphism F_i on some vector spaces V_i over some field of characteristic 0. One can try to compute $\zeta(X, T)$ using this formula.

Unfortunately, *étale cohomology* is usually¹ very hard to handle from an explicit computational point of view.

Much more progress has been made using *p*-adic cohomology in the spirit of Dwork. While there are several related approaches, we will focus on the construction of *Monsky-Washnitzer*, which has a close relationship with algebraic de Rham cohomology.

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Let U be a smooth affine variety over \mathbb{Q} (for simplicity). For some N > 0, it extends to a nice scheme² \mathfrak{U} over $\mathbb{Z}[1/N]$.

The algebraic de Rham cohomology $H^i_{dR}(U, \mathbb{Q})$ is the cohomology of (the global sections of) the de Rham complex

$$0 \to \mathcal{O}_U \to \Omega^1_{U/\mathbb{Q}} \to \Omega^2_{U/\mathbb{Q}} \to \dots \to \Omega^d_{U/\mathbb{Q}} \to 0 \qquad (d = \dim(U)).$$

Typically, the most interesting group occurs for i = d, as the other ones can be described in terms of cohomology of lower-dimensional varieties (Lefschetz hyperplane theorem).

To use $H^d_{dR}(U, \mathbb{Q})$ to compute zeta functions, it is critical to have a good algorithm to *reduce pole orders at infinity* in cohomology. Physicists may have more experience with this than number theorists...

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de Rham cohomology and zeta functions

Let p be a prime not dividing N. Then \mathfrak{U}_p (the reduction mod p) has a Frobenius endomorphism which does not typically lift to \mathfrak{U} ; but it does lift to a p-adic completion of \mathfrak{U} . By formally applying this action to differentials, one obtains³ endomorphisms F_i on $H^i_{dR}(U, \mathbb{Q}_p)$.

Theorem (Monsky-Washnitzer)

We have

$$\zeta(\mathfrak{U}_p, T) = \prod_{i=0}^{d} \det(1 - p^i F_i^{-1} T, H^i_{\mathsf{dR}}(U, \mathbb{Q}_p))^{(-1)^{i+1}}$$

Using *p*-adic numerical arithmetic, one can compute enough coefficients of $\zeta(\mathfrak{U}_p, \mathcal{T})$ with enough *p*-adic accuracy to determine them uniquely.

³This part is not formal; it amounts to saying that the de Rham cohomology doesn't change when you take the completion. This fails for the usual *p*-adic completion, so the slightly smaller *weak completion* is needed instead.

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For example, let $P \in \mathbb{Q}[x]$ be a polynomial of degree 2g+1 with no repeated roots, and put

$$U = \operatorname{Spec} \mathbb{Q}[x, y, z]/(y^2 - P(x), yz - 1).$$

Then for \pm the eigenspaces of the involution $y \mapsto -y$,

$$H^{1}_{dR}(U,\mathbb{Q}) = H^{1}_{dR}(U,\mathbb{Q})^{+} \oplus H^{1}_{dR}(U,\mathbb{Q})^{-}$$
$$= \left(\bigoplus_{i=0}^{2g} \mathbb{Q}\frac{dx}{y^{2}}\right) \oplus \left(\bigoplus_{i=0}^{2g-1} \mathbb{Q}\frac{dx}{y}\right)$$

It is easy to produce relations in $H^1_{dR}(U, \mathbb{Q})$ to reduce pole orders, e.g.,

$$\frac{(A(x)P(x) + B(x)P'(x))dx}{y^{2n+1}} \equiv \frac{A(x)dx}{y^{2n-1}} + \frac{2B'(x)dx}{(2n-1)y^{2n-1}}$$
$$0 \equiv \frac{(2dx^{d-1}P(x) + x^dP'(x))dx}{(2d+2g+1)y} = \frac{(x^{d+2g} + \cdots)dx}{y}.$$

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It is easy to produce relations in $H^1_{dR}(U,\mathbb{Q})$ to reduce pole orders, e.g.,

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For example, let $P \in \mathbb{Q}[x]$ be a polynomial of degree 2g+1 with no repeated roots, and put

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One can then (up to a suitably small *p*-adic error, in order to make the computation finite and tractable) apply this Frobenius lift formally to each basis differential of $H^i_{dR}(U, \mathbb{Q})$, reduce pole orders to get the matrix of action of Frobenius, and read off $\zeta(U, T)$. This is extremely efficient in practice! (Implementations exist in MAGMA, SAGE.)

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It is convenient to write differentials as degree 0 multiples of

$$\Omega = \sum_{i=0}^{n} (-1)^{i} x_{i} \, dx_{0} \wedge \cdots \wedge \widehat{dx_{i}} \wedge \cdots \wedge dx_{n}.$$

The reduction of pole orders is due to Griffiths-Dwork, using relations

$$\frac{\partial_i g}{f^m} \Omega \cong m \frac{g(\partial_i f)}{f^{m+1}} \Omega, \qquad \partial_i = \frac{\partial}{\partial x_i}.$$

In particular, one can construct a basis of $H^n_{dR}(U, \mathbb{Q})$ consisting of forms of the shape $g\Omega/f^i$ for $i \in \{1, \ldots, n\}$.

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One can then proceed as for hyperelliptic curves. However, a crude implementation of this idea (as in AKR, 2009) only yields practical results in a few small cases (e.g., n = 3, d = 4, p < 20).

A much better organization of the reduction process has been proposed by Harvey. Using this idea, Costa was able to calculate examples with n = 3, d = 4, $p < 2^{16}$.

Using similar methods, Costa also succeeded in computing a few examples with n = 4, d = 5, p < 20 in the Dwork pencil of quintic threefolds.

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Toric varieties

Let Δ be a polytope with vertices in \mathbb{Q}^n and choose a lattice $L \subset \mathbb{Q}^n$. We may then define a (projective, normal) toric variety X carrying an ample line bundle $\mathcal{O}(1)$ such that for $n = 0, 1, \dots, \Gamma(X, \mathcal{O}(n))$ is the Q-span of $L \cap n\Delta$. In particular, X contains an open dense torus⁴ $T = \operatorname{Spec} \mathbb{Q}[L]$ which acts on X, and X admits a locally closed stratification by T-orbits.

For example, if Δ is the standard unit simplex (whose vertices are 0 and the standard basis) and $L = \mathbb{Z}^n$, then $X = \mathbb{P}^n$ with the usual $\mathcal{O}(1)$.

More generally, for any positive integers a_0, \ldots, a_n , if we keep the same Δ but take $L = (a_1/a_0)\mathbb{Z} + \cdots + (a_n/a_0)\mathbb{Z}$, we get the weighted projective space $\mathbb{P}^n(a_0,\ldots,a_n)$.

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Recent progress in computing zeta functions

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Examples of interesting toric hypersurfaces

By a result of Reid, there are 95 different polytopes Δ for which the generic section of $\mathcal{O}(1)$ in X is a K3 surface. This list can also be found in Yonamura [Tôhoku, 1990] or this online database (with codimension=1):

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For f nondegenerate and U = T - Z, there is a simple analogue⁵ of the Griffiths-Dwork reduction process for computing $H^n_{dR}(U, \mathbb{Q})$: for any $\mathbb{Q}[L] \cong \mathbb{Q}[x_1^{\pm}, \ldots, x_n^{\pm}]$, $\Omega^n_{U/\mathbb{Q}}$ is generated by

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Controlled reduction for nondegenerate hypersurfaces

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By doing some extra linear algebra, we can construct extra relations which can be used to perform *controlled reduction*, in which sparseness is preserved: any form supported on a translate of $n\Delta$ reduces to another such form.

This is inspired by some related work of Harvey, who in the case of hyperelliptic curves obtained additional optimizations for the problem of computing L-functions (i.e., starting with a given U and reducing modulo many different primes). These optimizations are possible here too.

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